CPS506 - Comparative Programming Languages Implementation

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Why Managed Languages?

- productivity focus on the problem
- expressive languages functional, OO, declarative
- safety hard to get low-level details right

What Managed Languages?

- memory management usually garbage collected
- higher-level abstractions
- often interpreted/JIT
- often VM JavaVM and CLR are most well known

Which Managed Languages?

- OO Smalltalk, Java, Python, Ruby, C#, Scala, Javascript
- functional Elixir/Erlang, Haskell, SML, Ocaml, Racket/Scheme/LISP/Clojure
- array APL/J, R, MATLAB, Maple
- logic/declarative Prolog
- procedural/systems Go, Nim, Lua
- easier question what's not? C, C++, Rust, Zig, Odin, Jai

Modern Execution Structure?

• most machine architecture: PC, SP, other registers, memory

Dynamically Typed Languages

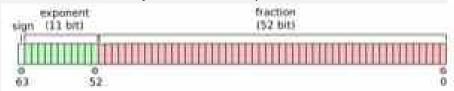
- values are typed
- some form of polymorphism parametric or OO
- means everything must be same size ideally a register
- could heap-allocate everything really bad for integers
- otherwise, need to tag values with type
- some hardware has had tags SPARC, B7700

Conventional tag

- modern architectures are byte-addressable
- heap objects will always be aligned say 8-byte boundaries
- can put tag in low bits, have integers shifted
- keep floating point values boxed

IEEE-FP tag

- modern processors have 64-bit integers, 64-bit pointers, and 64-bit IEEE floats
- IEEE floats have many Nan values exp all 1s 2⁵³ Nan values



- several ways to do NaN tagging/encoding
 - you can choose integers, pointers, or doubles to be naturally encoded
 - all the others be encoded with some shifting/adding
 - while integers and pointers are probably more common in most Smalltalk images
 - leaving doubles as naturally encoded means that FPU, vector instructions and/or GPUs can act directly on memory

IEEE-FP tag...

AST Smalltalk uses the following encoding based on the Sign+Exponent and Fraction bits:

S+E	F	F	F	Type
0000	0000	0000	0000	double +0
0000-7FEF	XXXX	XXXX	XXXX	double (positive)
7FF0	0000	0000	0000	+inf
7FF0-F	XXXX	XXXX	XXXX	NaN (unused)
8000	0000	0000	0000	double -0
8000-FFEF	XXXX	XXXX	XXXX	double (negative)
FFF0	0000	0000	0000	-inf
FFF2-F	XXXX	XXXX	xxxt	tagged literals
FFF2/3	XXXX	XXXX	XXXX	heap object
FFF4	0000	0001	0000	False
FFF6	0000	0010	0001	True
FFF8	1000	0000	0000	UndefinedObject
FFFA/B	XXXX	XXXX	XXXX	Symbol
FFFC/D	xxxx	XXXX	XXXX	Character
FFFE/F	XXXX	XXXX	XXXX	SmallInteger

Heaps

- sequential allocation very cheap but run out of space eventually
 - can work if we can compact out the freed memory eventually
- block allocation data structure to remember freed memory
 - many algorithms
 - external fragmentation free space that can't be allocated
 - internal fragmentation allocations larger than the object
- page allocation "pages" of uniform types
 - external fragmentation just at end of "page"
 - internal fragmentation should be none apart from alignment

Heap allocation

- manual allocation explicit malloc/free very error prone
- reference counting problem with cyclic structures, cascading-free
- non-moving collection
- compacting collector enables sequential allocation
- copying collector enables sequential allocation
- generational collection

Mark+Sweep Garbage Collection

- 2 phases
- mark phase go from roots to find all accessible data
- go through all object putting inaccessible into "free list"
- can be written to be mostly parallel
- can be conservative
- does not support sequential allocation
- significant fragmentation can exist
- allocation can be slow finding appropriate free space

Compacting Garbage Collection

- similar to mark+sweep with extra overhead to manage compacting
- sequential collector
- consolidate free space to prevent fragmentation and support sequential allocation

Generational Collection

- can be best of all worlds
- per-thread copying collector nursery + intermediate
- shared mark+sweep collector can be parallel

Copying Garbage Collection

- consolidate free space to prevent fragmentation and support sequential allocation
- sequential collector
- from roots collect all live objects into new area
- leaving "forwarding pointers" behind
- make the new space the current space
- only touches live data