CPS506 - Comparative Programming Languages Implementation

Dr. Dave Mason Department of Computer Science Ryerson University

©2022 Dave Mason © BY NO SA



Why Managed Languages?

- productivity focus on the problem
- expressive languages functional, OO, declarative
- safety hard to get low-level details right

What Managed Languages?

- memory management usually garbage collected
- higher-level abstractions
- often interpreted/JIT
- often VM JavaVM and CLR are most well known

Which Managed Languages?

- OO Smalltalk, Java, Python, Ruby, C#, Scala, Javascript
- functional Elixir/Erlang, Haskell, SML, Ocaml, Racket/Scheme/LISP/Clojure
- array APL/J, R, MATLAB, Maple
- logic/declarative Prolog
- procedural/systems Go, Nim, Lua
- easier question what's not? C, C++, Rust, Zig, Odin, Jai

Modern Execution Structure?

• most machine architecture: PC, SP, other registers, memory

Dynamically Typed Languages

- values are typed
- some form of polymorphism parametric or OO
- means everything must be same size ideally a register
- could heap-allocate everything really bad for integers
- otherwise, need to tag values with type
- some hardware has had tags SPARC, B7700

Conventional tag

- modern architectures are byte-addressable
- heap objects will always be aligned say 8-byte boundaries
- can put tag in low bits, have integers shifted
- keep floating point values boxed

IEEE-FP tag

- modern processors have 64-bit integers, 64-bit pointers, and 64-bit IEEE floats
- IEEE floats have many Nan values exp all 1s 2⁵³ Nan values



- several ways to do NaN tagging/encoding
 - you can choose integers, pointers, or doubles to be naturally encoded
 - all the others be encoded with some shifting/adding
 - while integers and pointers are probably more common in most Smalltalk images
 - leaving doubles as naturally encoded means that FPU, vector instructions and/or GPUs can act directly on memory

in a second and the second second of the eight _ ponone and .			
F	F	F	Туре
0000	0000	0000	double +0
XXXX	XXXX	XXXX	double (positive)
0000	0000	0000	+inf
XXXX	XXXX	xxxx	NaN (unused)
0000	0000	0000	double -0
XXXX	XXXX	xxxx	double (negative)
0000	0000	0000	-inf
XXXX	xxxx	xxxt	tagged literals
XXXX	XXXX	XXXX	heap object
0000	0001	0000	False
0000	0010	0001	True
1000	0000	0000	UndefinedObject
XXXX	XXXX	xxxx	Symbol
XXXX	XXXX	XXXX	Character
XXXX	XXXX	XXXX	SmallInteger
	F 0000 xxxx 0000 xxxx 0000 xxxx 0000 xxxx xxx 0000 0000 1000 xxxx xxxx xxxx	F F 0000 0000 xxxx xxxx xxxx xxxx 0000 0010 1000 0000 xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx	F F F 0000 0000 0000 xxxx xxxx xxxx 0000 0001 0000 0000 0010 0001 0000 0010 0001 1000 0000 0000 xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx xxxx

AST Smalltalk uses the following encoding based on the Sign+Exponent and Fraction bits:

Heaps

- sequential allocation very cheap but run out of space eventually
 - can work if we can compact out the freed memory eventually
- block allocation data structure to remember freed memory
 - many algorithms
 - external fragmentation free space that can't be allocated
 - internal fragmentation allocations larger than the object
- page allocation "pages" of uniform types
 - external fragmentation just at end of "page"
 - internal fragmentation should be none apart from alignment

Heap allocation

- manual allocation explicit malloc/free very error prone
- reference counting problem with cyclic structures, cascading-free
- non-moving collection
- compacting collector enables sequential allocation
- copying collector enables sequential allocation
- generational collection

Mark+Sweep Garbage Collection

- 2 phases
- mark phase go from roots to find all accessible data
- go through all object putting inaccessible into "free list"
- can be written to be mostly parallel
- can be conservative
- does not support sequential allocation
- significant fragmentation can exist
- allocation can be slow finding appropriate free space

Compacting Garbage Collection

- similar to mark+sweep with extra overhead to manage compacting
- sequential collector
- consolidate free space to prevent fragmentation and support sequential allocation

Copying Garbage Collection

- consolidate free space to prevent fragmentation and support sequential allocation
- sequential collector
- from roots collect all live objects into new area
- leaving "forwarding pointers" behind
- make the new space the current space
- only touches live data

Generational Collection

- can be best of all worlds
- per-thread copying collector nursery + intermediate
- shared mark+sweep collector can be parallel