

# CPS506 - Comparative Programming Languages

## Safety & Rust

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## Why Program Safety?

- economic costs
- health and safety
- personal information
- exploits and compromises

## How Safety?

- testing can only go so far
- dynamically typed languages are “safe”
- but static typing *can* provide more confidence
- ultimate is proof of correctness

## How Safe?

- static language safety dependent on type system
- “C” and “C++” are statically typed, but ...
  - null pointer exceptions
  - memory corruption
  - buffer overflows
- want expressive power too
- sometimes need all the performance you can get

## Rust

- need a safe systems-programming language
- minimal, predictable overhead
- strongly, statically typed
- no “undefined behaviour” a la C or C++ specs
- no data races

## History

- created at Mozilla - started 2010
- version 1.0 in May 2015
- Servo - browser currently at 100,000 LOC
- Dropbox internals
- Redox OS

## Paradigm

- imperative
- safe side-effects - even for multi-threading
- no mutable aliasing
- expressive type system
- no accidental run-time costs

## Syntax Rules

### 1 literals

- numbers: (un)signed ints, floats `-17`  
`3.141592`  
`13i8`
- characters: `'a'` - Unicode
- strings: `"this isn't \"hard\"!"`  
`r###"raw string with ' \ " # " ###"###"`  
`r#"useful for html <a href="fofof" & etc.."#"`
- arrays: `[1,2,3]` `[0;20]`
- slice: part of an array - `&a[..]` `&a[3..6]`
- tuples: `(1, "abc")`
- blocks/closures/lambdas:  
`|| 3`  
`|arg| arg-delta`

### 2 names

- upper/lower case, digits, underscore; case sensitive
- arguments to methods and blocks
- default immutable - `mut` keyword if needed
- snake-case for variables/functions/parameters
- camel-case for enum/struct/trait

### 3 functions

`fn name ( p1: t1, ... ) -> tr ( expr )`

## Statements - Conditionals

- **if**

```
let mut x = 3;
if x == 5 {
    x = 10
}
else {
    x += 1
}
```
- **match**

```
let x = 5;

match x {
    1 => println!("one"),
    2 => println!("two"),
    3 => println!("three"),
    4 => println!("four"),
    5 => println!("five"),
    _ => println!("something else"),
}
```

```
enum Message {
```

## Statements - Iterators

- **iterators - inlined (ranges, vectors, etc.)**

```
for (index, value) in (5..10).enumerate() {
    println!("index = {} and value = {}", index, value)
}

let lines = "hello\nworld".lines();

for (linenumber, line) in lines.enumerate() {
    println!("{}", linenumber, line);
}
```

## Statements - Loops

- **while**

```
let mut x = 5; // mut x: i32
let mut done = false; // mut done: bool

while !done {
    x += x - 3;

    println!("{}", x);

    if x % 5 == 0 {
        done = true;
    }
}
```
- **loop**

```
let mut x = 5; // mut x: i32

loop {
    x += x - 3;

    println!("{}", x);
}
```

## Structs

- **data containers**

```
struct MyData {
}

impl MyData {
}

inu
inu
```

## Traits

- like Haskell type-classes
- like Java interfaces - except not part of definition of base class
- trait *traitname*
- impl for any types

```
trait Shape {  
    fn draw(&self, Surface);  
    fn bounding_box(&self) -> BoundingBox;  
}  
impl Shape for i32 {  
  
}
```

## Trait Objects

```
struct Shape { ... }  
  
impl Shape {  
    fn draw(&self, u32) { ... }  
    fn bounding_box(&self) -> BoundingBox { ... }  
    fn default() -> &Self {  
    }  
}
```

```
let s = Shape{}  
s.draw(42)
```

- no hierarchy
- like Haskell type-classes

## Memory Safety

- no null pointers
  - way to create null pointers
  - `Option` enumerated type
- no dangling pointers
  - value lifetimes are calculated
  - *Rule 1: Every value has a single owner at any given time. You can move a value from one owner to another, but when a value's owner goes away, the value is freed along with it.*
  - *Rule 2: You can borrow a reference to a value, so long as the reference doesn't outlive the value (or equivalently, its owner).* Borrowed references are temporary pointers; they allow you to operate on values you don't own.
  - *Rule 3: You can only modify a value when you have exclusive access to it.*
- no memory leaks
  - value lifetimes are calculated
  - values freed when leave scope
  - additionally, reference-counted values
- no buffer overruns
  - no pointer arithmetic
  - slices used for partial arrays
  - arrays and slices are bounds checked

## Data Lifetimes

- small data implements Copy trait
- all arrays where element implements Copy trait
- everything else is moved
- assignment, parameter, result

## Lifetime...

```
fn make_vec() -> Vec<i32> {
    let mut vec = Vec::new();
    vec.push(0);
    vec.push(1);
    // scope ends, `vec` is destroyed
    vec // transfer ownership to the caller
}
fn print_vec(vec: &Vec<i32>) -> Vec<i32> {
    // the `vec` parameter is part of this scope, hence owned by `print_vec`
    // the `vec` parameter is borrowed for this scope
    for i in vec.iter() {
        println!("{}", i)
    }
    // now, `vec` is deallocated
    vec // now, pass ownership back
    // now, borrow ends
}
fn use_vec() {
    let vec = make_vec(); // take ownership of the vector
    let vec = print_vec(&vec); // pass ownershiplend access to `print_vec`
    // returned value is destroyed, as not used subsequently
    for i in vec.iter() { // Erroneously continue usinguse returned `vec`

        println!("{}", i * 2)
    }
    // scope ends, `vec` is destroyed
}
```

## Pointers

- Box<T> - heap allocated, moved
- &T and &mut T - references
- \*const T and \*mut T - C-like references - **unsafe**
- Rc - heap allocated immutable, clonable
- not sendable

## Lifetimes ...

- formally lifetimes of results are functions of parameter lifetimes
- fn bar<'a>(x: &'a i32) -> &'a i32
- lifetimes can sometimes be elided

```
struct Foo<'a> {
    x: &'a i32,
}

fn main() {
    let y = &5; // same as `let _y = 5; let y = &_y;`
    let f = Foo { x: y };

    println!("{}", f.x);
}

struct Foo<'a> {
    x: &'a i32,
}

impl<'a> Foo<'a> {
    fn x(&self) -> &'a i32 { self.x }
}
```

## Cells

- Cell<T> - mutable copy values
- RefCell<T> - mutable non-copy values
- usually used inside structs
- removes some of the simultaneous update guarantees
- not sendable

## Closures or Lambdas

- capture context

```
fn ten_times<F>(f: F) where F: Fn(i32) {
    for index in 0..10 {
        f(index);
    }
}

let greeting = "hello";
ten_times(|j| println!("{}", j), greeting, j);
```

## Modules

- mod *modname* ;
- mod *modname* { ... }

## Package Manager

- cargo
- creates for library or executable

## Macros

- hygienic, matching
  - zero or more items,
  - zero or more methods,
  - an expression,
  - a statement, or
  - a pattern.

```
let x: Vec<u32> = vec![1, 2, 3];
let x: Vec<u32> = {
    let mut temp_vec = Vec::new();
    temp_vec.push(1);
    temp_vec.push(2);
    temp_vec.push(3);
    temp_vec
};

macro_rules! vec {
    ( $( $x:expr ),* ) => {
        {
            let mut temp_vec = Vec::new();
            $(
                temp_vec.push($x);
            )
        }
    }
}
```

## Synchronous Types

- `Arc<T>` - heap allocated, clonable, sendable
- `Mutex<T>` - heap allocated, locked
- `RwLock<T>` - heap allocated, locked - read lock (multiple)

## Multi-processing

- channels
- mutex
- condition variables
- only types implementing `Send` can be sent or put in a `Mutex`
- means type system prevents data races

## Unsafe blocks

- occasionally need to reach under the covers
- including building the Rust library
- module or block can be declared unsafe to bypass type system

## Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- simple heap manager (no tracing or GC)
- array/slice bounds checking

## Evaluation

- Simplicity
  - Size of the grammar
  - Type system
  - complexity of navigating modules/classes
- Orthogonality
  - number of special syntax forms
  - number of special datatypes
- Extensibility
  - functional
  - syntactically
  - defining literals
  - overloading

## Zig

- another safe systems-programming language (also Odin, D, Nim, Jae)
- minimal, predictable overhead - even more than Rust
- statically typed, including array sizes
- no “undefined behaviour” a la C or C++ specs
- casts without `unsafe`
- 4 compilation models - Debug, ReleaseSafe, ReleaseSmall, ReleaseFast

## History

- created by Andrew Kelly
- version 0.10 in March 2022
- Zig Foundation funding development of self-hosting 1.0

## Paradigm

- imperative
- first-class types
- compile-time interpreter
- no accidental run-time costs
- no allocation without passing an allocator
- uses LLVM - dozens of targets (including wasm)



## Syntax Rules

### 1 literals

- numbers: (un)signed ints, floats `-17`  
`3.141592` (comptime - no default size) `@as(i56, 42)`
- characters: `'a'`
- UTF-8 strings are u8 arrays: `"this isn't \"hard\"!"`
- arrays: `[1, 2, 3]` `[5]u8{'h', 'e', 'l', 'l', 'o'}`  
`[_]u8{'w', 'o', 'r', 'l', 'd'}`
- slice: part of an array - `a[0..]` `a[3..6]`
- compile-time tuples (anonymous structs): `.{1, "abc"}`

### 2 names

- upper/lower case, digits, underscore; case sensitive  
`@"any thing!"`
- arguments to methods and blocks
- declarations: `const` or `var` - must be initialized (even if undefined)
- all variables must be used (even if `_ = variable`)
- snake\_case for variables/parameters
- camelCase for functions
- PascalCase enum/struct

### 3 functions

- `fn name ( p1 : t1 ... ) tr { expr }`
- `fn foo(x: i32) i32 { return x }`
- `var x: fn(i32) i32 = foo;`

## Statements - Conditionals

### • if

```
const expect = @import("std").testing.expect;
```

```
test "if statement" {
    const a = true;
    var x: u16 = 0;
    if (a) {
        x += 1;
    } else {
        x += 2;
    }
    try expect(x == 1);
}
```

### • switch

```
test "switch statement" {
    var x: i8 = 10;
    switch (x) {
        -1..1 => {
            x = -x;
        },
        10, 100 => {
            //special considerations must be made
            //when dividing signed integers
        }
    }
}
```

## Parsing

- `const` is used for types, errors, "normal values", modules
- modules are structs lazily imported from files/build-environment  
`const expect = @import("std").testing.expect;`
- values designated `pub` are visible to importers
- code is only cursorily parsed unless it is needed - very fast compile; lazy error detection; circular imports
- generics are done with type arguments-to/return-from functions
- no exceptions - errors or error-unions are return types for functions

## Statements - Loops

### • while

```
test "while with continue expression" {
    var sum: u8 = 0;
    var i: u8 = 1;
    while (i <= 10) : (i += 1) {
        sum += i;
    }
    try expect(sum == 55);
}
```

### • while with payload capture

```
var numbers_left: u32 = 4;
fn eventuallyNullSequence() ?u32 {
    if (numbers_left == 0) return null;
    numbers_left -= 1;
    return numbers_left;
}
```

```
test "while null capture" {
    var sum: u32 = 0;
    while (eventuallyNullSequence()) |value| {
        sum += value;
    }
    try expect(sum == 6); // 3 + 2 + 1
}
```

## Optional types and Iterators

- struct type with a next function with an optional in its return type
- returns null if no more values

```
const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
try expect(eql(u8, iter.next().?, "robust"));
try expect(eql(u8, iter.next().?, "optimal"));
try expect(eql(u8, iter.next().?, "reusable"));
try expect(eql(u8, iter.next().?, "maintainable"));
try expect(eql(u8, iter.next().?, ""));
try expect(iter.next() == null);

const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
var count : usize = 0;
while (iter.next()) |str| {
    count += str.len;
}
try expectEqual(count, 33);
```

## Structs

- data containers
- created by const declaration, or by function
- contain constants, variables, functions

```
const Point = struct {
    x: i32,
    y: i32,
    const Self = @This();
    pub fn new(x: i32, y: i32) Self {
        return Point{.x = x, .y = y};
    }
    pub fn abs(self: Self) Self {
        return new(if (self.x >= 0) self.x else -self.x,
                  if (self.y >= 0) self.y else -self.y);
    }
};
```

## Parametric Types

- functions can have types as parameters and can return types

```
pub fn Point_(comptime T: type) type {
    return struct {
        x: T,
        y: T,
        const Self = @This();
        pub fn new(x: T, y: T) Self {
            return .{.x = x, .y = y,};
        }
        pub fn abs(self: Self) Self {
            return new(if (self.x >= 0) self.x else -self.x,
                      if (self.y >= 0) self.y else -self.y);
        }
    };
}

test "parametric point" {
    const Point_i32 = Point_(i32);
    const p1 = Point_i32.new(3, -4);
    try expectEqual(p1.abs(), Point_i32.new(3, 4));
}
```

## Error handling

- no exceptions
- error returns
- must be handled - catch or try

## Memory Safety

- much weaker than Rust
- null pointers
  - but have to be recognized and dealt with
- dangling pointers
  - `defer` statement allows release adjacent to allocation
- buffer overruns
  - careful pointer arithmetic
  - slices used for partial arrays
  - arrays and slices are bounds-checked

## Modules

- just constant structs
- `@import("std") @import("heap.zig")`

## Macros

- no macros
- achieve similar ends with comptime first-class types

## Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- no automatic heap manager (no tracing or GC)
- array/slice bounds checking in safe/debug compilation modes
- undefined behaviour - detectable at compile time or run time

## Evaluation

- Simplicity
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- Extensibility
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