CPS506 - Comparative Programming Languages Safety & Rust

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Why Program Safety?

- economic costs
- health and safety
- personal information
- exploits and compromises

How Safety?

- testing can only go so far
- dynamically typed languages are "safe"
- but static typing can provide more confidence
- ultimate is proof of correctness

How Safe?

- static language safety dependent on type system
- "C" and "C++" are statically typed, but ...
 - null pointer exceptions
 - memory corruption
 - buffer overflows
- want expressive power too
- sometimes need all the performance you can get

Rust

- need a safe systems-programming language
- minimal, predictable overhead
- strongly, statically typed
- no "undefined behaviour" a la C or C++ specs
- no data races

History

- created at Mozilla started 2010
- version 1.0 in May 2015
- Servo browser currently at 100,000 LOC
- Dropbox internals
- Redox OS

Paradigm

- imperative
- safe side-effects even for multi-threading
- no mutable aliasing
- expressive type system
- no accidental run-time costs

Syntax Rules

literals

```
• numbers: (un)signed ints, floats -17 3.141592 13i8
```

characters: 'a' - Unicode

```
• strings: "this isn't \"hard\"!"
r###"raw string with ' \ " #" ##""###
r#"useful for html <a href="fofof" & etc.."#</pre>
```

• arrays: [1,2,3] [0;20]

- slice: part of an array &a[..] &a[3..6]
- tuples: (1, "abc")
- blocks/closures/lambdas:

```
|| 3
|arg| arg-delta
```

names

- upper/lower case, digits, underscore; case sensitive
- arguments to methods and blocks
- default immutable mut keyword if needed
- snake-case for variables/functions/parameters
- camel-case for enum/struct/trait
- functions

Statements - Conditionals

```
• if
 let mut x = 3;
 if x == 5 {
  x = 10
 else {
   x += 1
• match
 let x = 5;
 match x {
     1 => println!("one"),
     2 => println!("two"),
     3 => println!("three"),
     4 => println!("four"),
      5 => println!("five"),
     _ => println!("something else"),
 }
```

Statements - Loops

```
while
let mut x = 5; // mut x: i32
let mut done = false; // mut done: bool

while !done {
    x += x - 3;

    println!("{}", x);

    if x % 5 == 0 {
        done = true;
    }
}

loop
let mut x = 5; // mut x: i32

loop {
    x += x - 3;
}
```

Statements - Iterators

• iterators - inlined (ranges, vectors, etc.)

```
for (index, value) in (5..10).enumerate() {
    println!("index = {} and value = {}", index, value);
}
let lines = "hello\nworld".lines();

for (linenumber, line) in lines.enumerate() {
    println!("{}: {}", linenumber, line);
}
```

Structs

data containers

```
struct MyData {
}
impl MyData {
}
```

Traits

- like Haskell type-classes
- like Java interfaces except not part of definition of base class
- trait *traitname*
- impl for any types

```
trait Shape {
    fn draw(&self, Surface);
    fn bounding_box(&self) -> BoundingBox;
}
impl Shape for i32 {
}
```

Trait Objects

```
struct Shape { ... }
impl Shape {
    fn draw(&self, u32) { ... }
    fn bounding_box(&self) -> BoundingBox { ... }
    fn default() -> &Self {
}
let s = Shape{}
s.draw(42)
```

- no hierarchy
- like Haskell type-classes

Memory Safety

- no null pointers
 - way to create null pointers
 - Option enumerated type
- no dangling pointers
 - value lifetimes are calculated
 - Rule 1: Every value has a single owner at any given time. You can move a value from one owner to another, but when a value's owner goes away, the value is freed along with it.
 - Rule 2: You can borrow a reference to a value, so long as the reference doesn't outlive the value (or equivalently, its owner).
 Borrowed references are temporary pointers; they allow you to operate on values you don't own.
 - Rule 3: You can only modify a value when you have exclusive access to it.
- no memory leaks
 - value lifetimes are calculated
 - values freed when leave scope
 - additionally, reference-counted values
- no buffer overruns
 - no pointer arithmetic
 - slices used for partial arrays
 - arrays and slices are bounds-checked

Data Lifetimes

- small data implements Copy trait
- all arrays where element implements Copy trait
- everything else is moved
- assignment, parameter, result

Lifetime...

```
fn make_vec() -> Vec<i32> {
   let mut vec = Vec::new();
   vec.push(0);
   vec.push(1);
  // scope ends, 'vec' is destroyed
   vec // transfer ownership to the caller
fn print_vec(vec: &Vec<i32>) -> Vec<i32> {
// the 'vec' parameter is part of this scope, hence owned by 'print_vec'
// the 'vec' parameter is borrowed for this scope
   for i in vec.iter() {
       println!("{}", i)
       // now, 'vec' is deallocated
   vec // now, pass ownership back
       // now, borrow ends
fn use_vec() {
   let vec = make_vec(); // take ownership of the vector
   let vec = print_vec(&vec); // pass ownershiplend access to 'print_vec'
         // returned value is destroyed, as not used subsequently
   for i in vec.iter() { // Erroneously continue usinguse returned 'vec'
       println!("{}", i * 2)
   // scope ends, 'vec' is destroyed
}
```

Lifetimes ...

- formally lifetimes of results are functions of parameter lifetimes
- fn bar<'a>(x: &'a i32) -> &'a i32
- lifetimes can sometimes be elided

```
struct Foo<'a> {
    x: &'a i32,
}

fn main() {
    let y = &5; // same as 'let _y = 5; let y = &_y; '
    let f = Foo { x: y };

    println!("{}", f.x);
}

struct Foo<'a> {
    x: &'a i32,
}

impl<'a> Foo<'a> {
```

Pointers

- Box<T> heap allocated, moved
- &T and &mut T references
- *const T and *mut T C-like references unsafe
- Rc heap allocated immutable, clonable
- not sendable

Cells

- Cell<T> mutable copy values
- RefCell<T> mutable non-copy values
- usually used inside structs
- removes some of the simultaneous update guarantees
- not sendable

Closures or Lambdas

capture context

```
fn ten_times<F>(f: F) where F: Fn(i32) {
    for index in 0..10 {
        f(index);
    }
}
let greeting = "hello";
ten_times(|j| println!("{}, {}", greeting, j));
```

Modules

- mod *modname*;
- mod modname { · · · }

Package Manager

- cargo
- creates for library or executable

Macros

```
hygenic, matching
     • zero or more items,

    zero or more methods,

     • an expression,
     • a statement, or
      a pattern.
let x: Vec<u32> = vec![1, 2, 3];
let x: Vec<u32> = {
    let mut temp_vec = Vec::new();
    temp_vec.push(1);
    temp_vec.push(2);
    temp_vec.push(3);
    temp_vec
};
macro_rules! vec {
     ( $( $x:expr ),* ) => {
         {
             let mut temp_vec = Vec::new();
```

Synchronous Types

- Arc<T> heap allocated, clonable, sendable
- Mutex<T> heap allocated, locked
- RwLock<T> heap allocated, locked read lock (multiple)

Multi-processing

- channels
- mutex
- condition variables
- only types implementing Send can be sent or put in a Mutex
- means type system prevents data races

Unsafe blocks

- occasionally need to reach under the covers
- including building the Rust library
- module or block can be declared unsafe to bypass type system

Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- simple heap manager (no tracing or GC)
- array/slice bounds checking

Evaluation

- Simplicity
 - Size of the grammar
 - Type system
 - complexity of navigating modules/classes
- Orthogonality
 - number of special syntax forms
 - number of special datatypes
- Extensibility
 - functional
 - syntactically
 - defining literals
 - overloading

Zig

- another safe systems-programming language (also Odin, D, Nim, Jae)
- minimal, predictable overhead even more than Rust
- statically typed, including array sizes
- no "undefined behaviour" a la C or C++ specs
- casts without unsafe
- 4 compilation models Debug, ReleaseSafe, ReleaseSmall, ReleaseFast

History

- created by Andrew Kelly
- version 0.10 in March 2022
- Zig Foundation funding development of self-hosting 1.0

Paradigm

- imperative
- first-class types
- compile-time interpreter
- no accidental run-time costs
- no allocation without passing an allocator
- uses LLVM dozens of targets (including wasm)

Syntax Rules

- literals
 - numbers: (un)signed ints, floats -17
 3.141592 (comptime no default size) @as (i56, 42)
 - characters: 'a'
 - UTF-8 strings are u8 arrays: "this isn't \"hard\"!"
 - arrays: [1,2,3] [5]u8{'h','e','l','l','o']
 [_]u8{'w','o','r','l','d']
 - slice: part of an array a [0..] a [3..6]
 - compile-time tuples (anonymous structs): . {1, "abc"}
- names
 - upper/lower case, digits, underscore; case sensitive @"any thing!"
 - arguments to methods and blocks
 - declarations: const or var must be initialized (even if undefined)
 - all variables must be used (even if _ = variable
 - snake_case for variables/parameters
 - camelCase for functions
 - PascalCase enum/struct
- functions
 - fn name (p1:t1...) tr { expr }
 - fn foo(x: i32) i32 { return x }

Parsing

- const is used for types, errors, "normal values", modules
- modules are structs lazily imported from files/build-environment const expect = @import("std").testing.expect;
- values designated pub are visible to importers
- code is only cursorily parsed unless it is needed very fast compile; lazy error detection; circular imports
- generics are done with type arguments-to/return-from functions
- no exceptions errors or error-unions are return types for functions

Statements - Conditionals

• if const expect = @import("std").testing.expect; test "if statement" { const a = true; var x: u16 = 0;if (a) { x += 1;} else { x += 2;try expect (x == 1); switch test "switch statement" { var x: i8 = 10;switch (x) { $-1...1 => {$ x = -x;10, 100 => { //special considerations must be made //when dividing signed integers x = @divExact(x, 10);

Statements - Loops

```
while
  test "while with continue expression" {
      var sum: u8 = 0;
      var i: u8 = 1;
      while (i \le 10): (i += 1) {
          sum += i;
      try expect (sum == 55);
while with payload capture
  var numbers_left: u32 = 4;
  fn eventuallyNullSequence() ?u32 {
      if (numbers_left == 0) return null;
      numbers_left -= 1;
      return numbers_left;
  }
  test "while null capture" {
      var sum: u32 = 0;
      while (eventuallyNullSequence()) |value| {
          sum += value;
      var c: :134 = /;
      trv_{expect(sum_{eq})}^{7} = 6); // 3 + 2 + 1
```

Optional types and Iterators

- struct type with a next function with an optional in its return type
- returns null if no more values

```
const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
try expect(eql(u8, iter.next().?, "robust"));
try expect(eql(u8, iter.next().?, "optimal"));
try expect(eql(u8, iter.next().?, "reusable"));
try expect(eql(u8, iter.next().?, "maintainable"));
try expect(eql(u8, iter.next().?, ""));
try expect(iter.next() == null);
const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
var count : usize = 0;
while (iter.next()) |str| {
   count += str.len;
}
try expectEqual(count, 33);
```

Structs

- data containers
- created by const declaration, or by function
- contain constants, variables, functions

Parametric Types

functions can have types as parameters and can return types

```
pub fn Point_(comptime T: type) type {
    return struct {
        x: T,
        y: T,
        const Self = @This();
        pub fn new(x: T, y: T) Self {
            return . \{.x = x, .y = y,\};
        pub fn abs(self: Self) Self {
            return new(if (self.x>=0) self.x else -self.x,
                        if (self.y>=0) self.y else -self.y);
    };
test "parametric point" {
    const Point_i32 = Point_(i32);
    const p1 = Point_i32.new(3, -4);
    try expectEqual(p1.abs(), Point_i32.new(3,4));
}
```

Error handling

- no exceptions
- error returns
- must be handled catch or try

Memory Safety

- much weaker than Rust
- null pointers
 - but have to be recognized and dealt with
- dangling pointers
 - defer statement allows release adjacent to allocation
- buffer overruns
 - careful pointer arithmetic
 - slices used for partial arrays
 - arrays and slices are bounds-checked

Modules

- just constant structs
- @import("std") @import("heap.zig")

Macros

- no macros
- achieve similar ends with comptime first-class types

Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- no automatic heap manager (no tracing or GC)
- array/slice bounds checking in safe/debug compilation modes
- undefined behaviour detectable at compile time or run time

Evaluation

- Simplicity
 - Size of the grammar
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