

CPS506 - Comparative Programming Languages

Safety & Rust

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Why Program Safety?

- economic costs
- health and safety
- personal information
- exploits and compromises

How Safety?

- testing can only go so far
- dynamically typed languages are “safe”
- but static typing *can* provide more confidence
- ultimate is proof of correctness

How Safe?

- static language safety dependent on type system
- “C” and “C++” are statically typed, but ...
 - null pointer exceptions
 - memory corruption
 - buffer overflows
- want expressive power too
- sometimes need all the performance you can get

- need a safe systems-programming language
- minimal, predictable overhead
- strongly, statically typed
- no “undefined behaviour” a la C or C++ specs
- no data races

History

- created at Mozilla - started 2010
- version 1.0 in May 2015
- Servo - browser currently at 100,000 LOC
- Dropbox internals
- Redox OS

Paradigm

- imperative
- safe side-effects - even for multi-threading
- no mutable aliasing
- expressive type system
- no accidental run-time costs

Syntax Rules

1 literals

- numbers: (un)signed ints, floats -17
3.141592
13i8
- characters: 'a' - Unicode
- strings: "this isn't \"hard\"!"
r###"raw string with ' \ " # " ##"###"
r#"useful for html <a href="fofof" & etc.."#"
- arrays: [1,2,3] [0;20]
- slice: part of an array - &a[..] &a[3..6]
- tuples: (1, "abc")
- blocks/closures/lambdas:
|| 3
|arg| arg-delta

2 names

- upper/lower case, digits, underscore; case sensitive
- arguments to methods and blocks
- default immutable - `mut` keyword if needed
- snake-case for variables/functions/parameters
- camel-case for enum/struct/trait

Statements - Conditionals

- if

```
let mut x = 3;
if x == 5 {
    x = 10
}
else {
    x += 1
}
```

- match

```
let x = 5;
```

```
match x {
    1 => println!("one"),
    2 => println!("two"),
    3 => println!("three"),
    4 => println!("four"),
    5 => println!("five"),
    _ => println!("something else")
}
```

Statements - Loops

- while

```
let mut x = 5; // mut x: i32
let mut done = false; // mut done: bool
```

```
while !done {
    x += x - 3;

    println!("{}", x);

    if x % 5 == 0 {
        done = true;
    }
}
```

- loop

```
let mut x = 5; // mut x: i32
```

```
loop {
    x += x - 3;
}
```

Statements - Iterators

- iterators - inlined (ranges, vectors, etc.)

```
for (index, value) in (5..10).enumerate() {  
    println!("index = {} and value = {}", index, value)  
}
```

```
let lines = "hello\nworld".lines();
```

```
for (linenumber, line) in lines.enumerate() {  
    println!("{}", line, linenumber);  
}
```

Structs

- data containers

```
struct MyData {  
  
}  
impl MyData {  
  
}
```

Traits

- like Haskell type-classes
- like Java interfaces - except not part of definition of base class
- `trait traitname`
- `impl` for any types

```
trait Shape {  
    fn draw(&self, Surface);  
    fn bounding_box(&self) -> BoundingBox;  
}  
impl Shape for i32 {  
  
}
```

Trait Objects

```
struct Shape { ... }
```

```
impl Shape {  
    fn draw(&self, u32) { ... }  
    fn bounding_box(&self) -> BoundingBox { ... }  
    fn default() -> &Self {  
    }  
}
```

```
let s = Shape{}  
s.draw(42)
```

- no hierarchy
- like Haskell type-classes

Memory Safety

- no null pointers
 - way to create null pointers
 - `Option` enumerated type
- no dangling pointers
 - value lifetimes are calculated
 - *Rule 1: Every value has a single owner at any given time. You can move a value from one owner to another, but when a value's owner goes away, the value is freed along with it.*
 - *Rule 2: You can borrow a reference to a value, so long as the reference doesn't outlive the value (or equivalently, its owner). Borrowed references are temporary pointers; they allow you to operate on values you don't own.*
 - *Rule 3: You can only modify a value when you have exclusive access to it.*
- no memory leaks
 - value lifetimes are calculated
 - values freed when leave scope
 - additionally, reference-counted values
- no buffer overruns
 - no pointer arithmetic

Data Lifetimes

- small data implements Copy trait
- all arrays where element implements Copy trait
- everything else is moved
- assignment, parameter, result

Lifetime...

```
fn make_vec() -> Vec<i32> {
    let mut vec = Vec::new();
    vec.push(0);
    vec.push(1);
    // scope ends, `vec` is destroyed
    vec // transfer ownership to the caller
}

fn print_vec(vec: &Vec<i32>) -> Vec<i32> {
    // the `vec` parameter is part of this scope, hence owned by `print_vec`
    // the `vec` parameter is borrowed for this scope
    for i in vec.iter() {
        println!("{}", i)
    }
    // now, `vec` is deallocated
    vec // now, pass ownership back
    // now, borrow ends
}

fn use_vec() {
    let vec = make_vec(); // take ownership of the vector
    let vec = print_vec(&vec); // pass ownership back
    // returned value is destroyed, as not used subsequently
    for i in vec.iter() { // Erroneously continue using use returned `vec`

        println!("{}", i * 2)
    }
}
```

Lifetimes ...

- formally lifetimes of results are functions of parameter lifetimes
- `fn bar<'a>(x: &'a i32) -> &'a i32`
- lifetimes can sometimes be elided

```
struct Foo<'a> {  
    x: &'a i32,  
}
```

```
fn main() {  
    let y = &5; // same as `let _y = 5; let y = &_y;`  
    let f = Foo { x: y };  
  
    println!("{}", f.x);  
}
```

```
struct Foo<'a> {  
    x: &'a i32,  
}
```

Pointers

- `Box<T>` - heap allocated, moved
- `&T` and `&mut T` - references
- `*const T` and `*mut T` - C-like references - **unsafe**
- `Rc` - heap allocated immutable, clonable
- not sendable

- `Cell<T>` - mutable copy values
- `RefCell<T>` - mutable non-copy values
- usually used inside structs
- removes some of the simultaneous update guarantees
- not sendable

Closures or Lambdas

- capture context

```
fn ten_times<F>(f: F) where F: Fn(i32) {  
    for index in 0..10 {  
        f(index);  
    }  
}
```

```
let greeting = "hello";  
ten_times(|j| println!("{}", {}, greeting, j));
```

Modules

- `mod modname ;`
- `mod modname { ... }`

Package Manager

- cargo
- creates for library or executable

Macros

- hygienic, matching
 - zero or more items,
 - zero or more methods,
 - an expression,
 - a statement, or
 - a pattern.

```
let x: Vec<u32> = vec![1, 2, 3];
```

```
let x: Vec<u32> = {  
    let mut temp_vec = Vec::new();  
    temp_vec.push(1);  
    temp_vec.push(2);  
    temp_vec.push(3);  
    temp_vec  
};
```

```
macro_rules! vec {  
    ( $( $x:expr ),* ) => {  
        {  
            let mut temp_vec = Vec::new();
```


Synchronous Types

- `Arc<T>` - heap allocated, clonable, sendable
- `Mutex<T>` - heap allocated, locked
- `RwLock<T>` - heap allocated, locked - read lock (multiple)

Multi-processing

- channels
- mutex
- condition variables
- only types implementing `Send` can be sent or put in a Mutex
- means type system prevents data races

Unsafe blocks

- occasionally need to reach under the covers
- including building the Rust library
- module or block can be declared unsafe to bypass type system

Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- simple heap manager (no tracing or GC)
- array/slice bounds checking

Evaluation

- Simplicity
 - Size of the grammar
 - Type system
 - complexity of navigating modules/classes
- Orthogonality
 - number of special syntax forms
 - number of special datatypes
- Extensibility
 - functional
 - syntactically
 - defining literals
 - overloading

- another safe systems-programming language (also Odin, D, Nim, Jae)
- minimal, predictable overhead - even more than Rust
- statically typed, including array sizes
- no “undefined behaviour” a la C or C++ specs
- casts without `unsafe`
- 4 compilation models - Debug, ReleaseSafe, ReleaseSmall, ReleaseFast

History

- created by Andrew Kelly
- version 0.10 in March 2022
- Zig Foundation funding development of self-hosting 1.0

Paradigm

- imperative
- first-class types
- compile-time interpreter
- no accidental run-time costs
- no allocation without passing an allocator
- uses LLVM - dozens of targets (including wasm)

Syntax Rules

1 literals

- numbers: (un)signed ints, floats -17
3.141592 (comptime - no default size) @as(i56, 42)
- characters: 'a'
- UTF-8 strings are u8 arrays: "this isn't \"hard\"!"
- arrays: [1, 2, 3] [5]u8{'h', 'e', 'l', 'l', 'o'}
[_]u8{'w', 'o', 'r', 'l', 'd'}
- slice: part of an array - a[0..] a[3..6]
- compile-time tuples (anonymous structs): .{1, "abc"}

2 names

- upper/lower case, digits, underscore; case sensitive
@"any thing!"
- arguments to methods and blocks
- declarations: `const` or `var` - must be initialized (even if undefined)
- all variables must be used (even if `_ = variable`)
- snake_case for variables/parameters
- camelCase for functions
- PascalCase enum/struct

3 functions

Parsing

- `const` is used for types, errors, “normal values”, modules
- modules are `structs` lazily imported from files/build-environment
`const expect = @import("std").testing.expect;`
- values designated `pub` are visible to importers
- code is only cursorily parsed unless it is needed - very fast compile; lazy error detection; circular imports
- generics are done with type arguments-to/return-from functions
- no exceptions - errors or error-unions are return types for functions

Statements - Conditionals

- if

```
const expect = @import("std").testing.expect;
```

```
test "if statement" {
    const a = true;
    var x: u16 = 0;
    if (a) {
        x += 1;
    } else {
        x += 2;
    }
    try expect(x == 1);
}
```

- switch

```
test "switch statement" {
    var x: i8 = 10;
    switch (x) {
        -1...1 => {
            x = -x;
        },
        10, 100 => {
            //special considerations must be made
        }
    }
}
```

Statements - Loops

- while

```
test "while with continue expression" {
  var sum: u8 = 0;
  var i: u8 = 1;
  while (i <= 10) : (i += 1) {
    sum += i;
  }
  try expect(sum == 55);
}
```

- while **with payload capture**

```
var numbers_left: u32 = 4;
fn eventuallyNullSequence() ?u32 {
  if (numbers_left == 0) return null;
  numbers_left -= 1;
  return numbers_left;
}

test "while null capture" {
  var sum: u32 = 0;
  while (eventuallyNullSequence()) |value| {
    sum += value;
  }
}
```

Optional types and Iterators

- struct type with a next function with an optional in its return type
- returns `null` if no more values

```
const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
try expect(eql(u8, iter.next().?, "robust"));
try expect(eql(u8, iter.next().?, "optimal"));
try expect(eql(u8, iter.next().?, "reusable"));
try expect(eql(u8, iter.next().?, "maintainable"));
try expect(eql(u8, iter.next().?, ""));
try expect(iter.next() == null);

const text = "robust, optimal, reusable, maintainable, ";
var iter = std.mem.split(u8, text, ", ");
var count : usize = 0;
while (iter.next()) |str| {
    count += str.len;
}
try expectEqual(count, 33);
```

Structs

- data containers
- created by `const` declaration, or by function
- contain constants, variables, functions

```
const Point = struct {
    x: i32,
    y: i32,
    const Self = @This();
    pub fn new(x: i32, y: i32) Self {
        return Point{.x = x, .y = y};
    }
    pub fn abs(self: Self) Self {
        return new(if (self.x>=0) self.x else -self.x,
                  if (self.y>=0) self.y else -self.y);
    }
};
```

Parametric Types

- functions can have types as parameters and can return types

```
pub fn Point_(comptime T: type) type {
  return struct {
    x: T,
    y: T,
    const Self = @This();
    pub fn new(x: T, y: T) Self {
      return .{.x = x, .y = y,};
    }
    pub fn abs(self: Self) Self {
      return new(if (self.x >= 0) self.x else -self.x,
                 if (self.y >= 0) self.y else -self.y);
    }
  };
}

test "parametric point" {
  const Point_i32 = Point_(i32);
  const p1 = Point_i32.new(3, -4);
  try expectEqual(p1.abs(), Point_i32.new(3, 4));
}
```

Error handling

- no exceptions
- error returns
- must be handled - `catch` or `try`

Memory Safety

- much weaker than Rust
- null pointers
 - but have to be recognized and dealt with
- dangling pointers
 - `defer` statement allows release adjacent to allocation
- buffer overruns
 - careful pointer arithmetic
 - slices used for partial arrays
 - arrays and slices are bounds-checked

Modules

- just constant structs
- `@import("std") @import("heap.zig")`

Macros

- no macros
- achieve similar ends with comptime first-class types

Pragmatics

- predictable, high performance
- almost no run-time system required
- native compilation
- no automatic heap manager (no tracing or GC)
- array/slice bounds checking in safe/debug compilation modes
- undefined behaviour - detectable at compile time or run time

Evaluation

- Simplicity
 - Size of the grammar
 - Type system
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